

Supplemental Notes

EE503 week 02

Dr. Franzke

HW #2

① HW #2 handout.

Argument assignment.

② Leon-Garcia

③ Gubner

Topics

- Rules of Inference (Mathematical argumentation)
- Borel Sigma-Algebra CIA.
- Conditional Probability
- Theorem of total probability and Bayes Theorem

Rules of Inference

(the logical schema)

modus Ponens

Premise 1: $P \rightarrow Q$

Premise 2: P

Conclusion: Q

Theorem: $[(P \rightarrow Q) \& P] \rightarrow Q$

Prf: by truth table

Modus Tollens

Premise 1: $P \rightarrow Q$

Premise 2: $\sim Q$

Conclusion: $\sim P$

Theorem: $[(P \rightarrow Q) \& \sim Q] \rightarrow \sim P$

Conditional: $P \rightarrow Q$

Contra positive: $\sim Q \rightarrow \sim P$

Converse: $Q \rightarrow P$

Theorem: "Material implication"

$$P \rightarrow Q = \sim P \vee Q$$

Defn: An argument is VALID iff the premises logically imply the conclusion

(can't make if-part TRUE and then-part FALSE)

Defn: An argument is SOUND iff it is valid and all premises are true

CIA GUT CAT

- mnemonic for generating sigma-algebras.

Thm: $\mathcal{Q}_\alpha \subset 2^X$ are S.A. $\longrightarrow \bigcap_\alpha \mathcal{Q}_\alpha$ is a S.A.

\therefore S.A.'s are closed under intersection

CIA: Pick any $\mathcal{Q} \subset 2^X$. Then \mathcal{Q} generates the S.A. $\sigma(\mathcal{Q})$:

$$\sigma(\mathcal{Q}) = \bigcap_\alpha \mathcal{Q}_\alpha, \text{ where}$$

C: Containment: $\mathcal{Q} \subset \mathcal{Q}_\alpha \quad \forall \alpha$

I: Intersection: $\bigcap_\alpha \mathcal{Q}_\alpha$

A: Sigma-algebras: \mathcal{Q}_α is a S.A. (GUT) $\forall \alpha$.

Defn: $\mathcal{B}(\mathbb{R})$ is the Borel S.A. on \mathbb{R} . iff

$$\mathcal{B}(\mathbb{R}) = \sigma\left(\{(-\infty, a] : a \in \mathbb{R}\}\right)$$

$\therefore \mathcal{B}(\mathbb{R})$ is the "interval S.A."

Fact: $|\mathcal{B}(\mathbb{R})| = \mathfrak{c} = |\mathbb{R}| < 2^{\mathfrak{c}} = |2^{\mathbb{R}}|$.

\uparrow
"power of the continuum"

$\therefore \mathcal{B}(\mathbb{R})$ is not "too-big" for a prob. space

(unlike $2^{\mathbb{R}}$)

Defn: $f: X \rightarrow Y$ is 1-to-1 ("Injective") iff
 $\forall x \forall z: f(x) = f(z) \rightarrow x = z.$

Defn: $f: X \rightarrow Y$ is ONTO ("Surjective") iff
 $\forall y \in Y \exists x \in X: y = f(x)$ ($f(X) = Y$)

Defn: $f: X \rightarrow Y$ is a BIJECTION ("1-to-1 correspondence")
iff f is 1-to-1 and onto.

$|A| =$ cardinality ("size") of set A .

$\mathbb{N} = \{1, 2, 3, \dots\}$ natural numbers

$\mathbb{Z} = \{0, \pm 1, \pm 2, \pm 3, \dots\}$ integers

$$\mathbb{Z}^+ = \mathbb{N}$$

$\mathbb{R} = (-\infty, \infty)$ reals

Defn: A is finite iff $A \overset{1\text{-to-1}}{\underset{\text{onto}}{\longleftrightarrow}} S$ for some $S \subset \mathbb{N}$
such that $A = \{a_1, \dots, a_n\}$ for some $n \in \mathbb{N}$ (or $A = \emptyset$)
Else A is infinite.

Fact: A infinite $\iff A \overset{1\text{-to-1}}{\underset{\text{onto}}{\longleftrightarrow}} B$ for some $B \subset A$ and $B \neq A$.

Cantor's theorem: $|X| < |2^X|$

Defn: A is denumerable iff $A \overset{1\text{-to-1}}{\underset{\text{onto}}{\longleftrightarrow}} \mathbb{N}$ (e.g. \mathbb{Z})

Defn: A is countable iff (1) A is finite, or
(2) A is denumerable

Facts: $|\mathbb{Z}| = |\mathbb{N}| = \aleph_0$ (aleph-nought)

$$|\mathbb{R}| = |\mathbb{Z}^{\mathbb{Z}}| = c = \aleph_1$$

↑ power of the continuum

$$|A| = \aleph_k \implies |A^A| = \aleph_{k+1} \quad (k=0,1,\dots)$$

Cantor's Continuum Hypothesis There is no ω such that

$$\aleph_k < \omega < \aleph_{k+1}$$

Fact: There is no S.A. \mathbb{Q} : $|\mathbb{Q}| = |\mathbb{Z}|$.

Ex: Cannot pick an integer "at random"

Defn: Suppose (Ω, \mathcal{A}, P) is a probability space.
 Events A and B are (statistically) independent
 w.r.t. P iff $P(A \cap B) = P(A) \cdot P(B)$

"joint factors into marginals"

Defn: Conditional Probability $P(B|A) = \frac{P(A \cap B)}{P(A)}$ if $P(A) > 0$.

$\therefore A$ & B independent $\iff P(B|A) = P(B)$
 $\iff P(A|B) = P(A)$

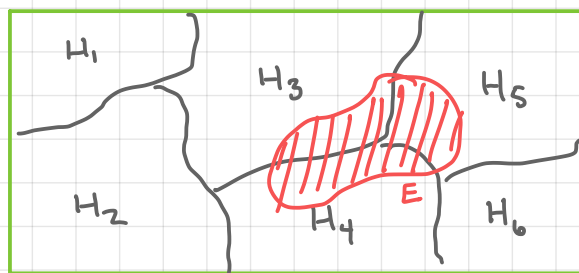
Defn: Suppose (Ω, \mathcal{A}) is a measurable space. Then
 $\{H_k\} \subset \mathcal{A}$ partition Ω ("is a partition") iff

(1) $\bigcup_k H_k = \Omega$

(2) pairwise-disjoint $H_i \cap H_j = \emptyset$ if $i \neq j$

Theorem: (of Total Probability)

$P(E) = \sum_k P(H_k) \cdot P(E|H_k)$ if $\{H_k\}$ partition Ω



x

Hypotheses: H_k

Evidence: E

Prf: $P(E) = P[E \cap \Omega]$
 $= P[E \cap (\bigcup_k H_k)]$ since $\{H_k\}$ partition Ω

$$= P\left(\bigcup_k (E \cap H_k)\right)$$

$$= \sum_k P(E \cap H_k) \quad \text{since } \{H_k\} \text{ partition } \Omega$$

$$= \sum_k P(H_k) \cdot P(E|H_k) \quad \text{defn of cond. prob.}$$

QED

$$P(B|A) = \frac{P(A \cap B)}{P(A)}$$

Later in this course:

Total Expectation

$$E_x[X] = E_y[E[X|Y]]$$

Total Variance

$$V_x[X] = E_y[V[X|Y]] + V_y[E[X|Y]].$$

extremely important for problem solving

Special case: $\Omega = A \cup A^c$

$$\therefore P(B) \stackrel{\text{total prob.}}{=} P(A) \cdot P(B|A) + P(A^c) \cdot P(B|A^c)$$

Note: for any B.

Theorem: (Bayes' Theorem)

$$P(H_j|E) = \frac{P(E|H_j) \cdot P(H_j)}{\sum_k P(E|H_k) \cdot P(H_k)} \quad \text{if } \{H_k\} \text{ partitions } \Omega.$$

Prf: $P(H_j|E) = \frac{P(E \cap H_j)}{P(E)}$ defn of cond. prob

$$= \frac{P(H_j) \cdot P(E|H_j)}{P(E)} \quad \text{" " "}$$

$$= \frac{P(H_j) \cdot P(E|H_j)}{\sum_k P(H_k) \cdot P(E|H_k)} \quad \text{total prob.}$$

QED

$P(H_j)$: Prior

$P(H_j|E)$: Posterior

$P(E|H_j)$: Likelihood

$$\therefore P(H|E) = \frac{a}{a+b}$$

$$P(H^c|E) = \frac{b}{a+b}$$

Mnemonic:

"Party Unconditionally To Conquer Bayes"

P: Partition $\{H_k\}$ or $A \cup A^c$

U: Unconditional probability $P(B)$

T: Total Probability

C: Conditional Probability $P(B|A)$

B: Bayes Theorem

Issue Spotting Sequence

- Is there a partition?

- Unconditional probability $P(B)$?

\therefore Use total probability

- Conditional probability $P(B|A)$?

\therefore Use Bayes Theorem

Theorem: ("Booles inequality") $P\left(\bigcup_{k=1}^n A_k\right) \leq \sum_{k=1}^n P(A_k)$

Theorem: ("Multiplication theorem")

$$P\left(\bigcap_{k=1}^n A_k\right) = P(A_1) \cdot P(A_2 | A_1) \cdots P(A_n | A_1 \cap \cdots \cap A_{n-1})$$

Prf: (by induction on $n = 2, 3, 4, \dots$)

Basis Step: $n = 2$.

$$P(A_1 \cap A_2) = P(A_1) \cdot P(A_2 | A_1) \quad \text{by defn cond. prob.}$$

QED-basis

Induction Step: - Assume holds for n .

- Derive that it holds for $n+1$

Induction Hypothesis (IH):

$$P\left(\bigcap_{k=1}^n A_k\right) = P(A_1) \cdot P(A_2 | A_1) \cdots P(A_n | A_1 \cap \cdots \cap A_{n-1})$$

$$\begin{aligned} \Rightarrow P\left(\bigcap_{k=1}^{n+1} A_k\right) &= P\left(\bigcap_{k=1}^n A_k \cap A_{n+1}\right) \\ &\stackrel{\text{assoc.}}{=} P\left(\left(\bigcap_{k=1}^n A_k\right) \cap A_{n+1}\right) \\ &= P\left(\bigcap_{k=1}^n A_k\right) \cdot P\left(A_{n+1} \mid \bigcap_{k=1}^n A_k\right) \\ &= P(A_1) \cdot P(A_2 | A_1) \cdots P(A_n | A_1 \cap \cdots \cap A_{n-1}) \\ &\quad \cdot P\left(A_{n+1} \mid \bigcap_{k=1}^n A_k\right) \end{aligned}$$

QED.

Odds form of Bayes' Theorem:

$$\frac{P(H|E)}{P(H^c|E)} = \frac{\overbrace{P(E|H)}^{\text{likelihood ratio}}}{P(E|H^c)} \cdot \frac{P(H)}{P(H^c)}$$

$$\text{or } O(H|E) = \underbrace{L(H|H^c)}_{\text{likelihood ratio}} \cdot O(H)$$

$$\text{where } O(A) = \text{Odds}(A) = \frac{P(A)}{P(A^c)} \quad \therefore P(A) = \frac{O(A)}{1+O(A)}$$

Two-class Classification: H vs. H^c

Then can write the Bayes' Theorem posterior $P(H|E)$ in terms of the odds form Bayes' theorem term:

$$O(H|E) = \frac{P(H|E)}{P(H^c|E)} = \frac{P(E|H) \cdot P(H)}{P(E|H^c) \cdot P(H^c)}$$

Theorem: The posterior $P(H|E)$ has the logistic form.

$$P(H|E) = \frac{1}{1 + e^{-\phi(E)}} \quad \text{where } \phi(E) = \ln O(H|E)$$

Prf:

$$P(H|E) = \frac{P(H) \cdot P(E|H)}{P(H) \cdot P(E|H) + P(H^c) \cdot P(E|H^c)}$$

$$= \frac{1}{1 + \frac{P(H^c) \cdot P(E|H^c)}{P(H) \cdot P(E|H)}}$$

$$= \frac{1}{1 + \frac{1}{O(H|E)}}$$

$$= \frac{1}{1 + e^{\ln(1/\sigma(H|\epsilon))}}$$

$$= \frac{1}{1 + e^{-\ln \sigma(H|\epsilon)}}$$

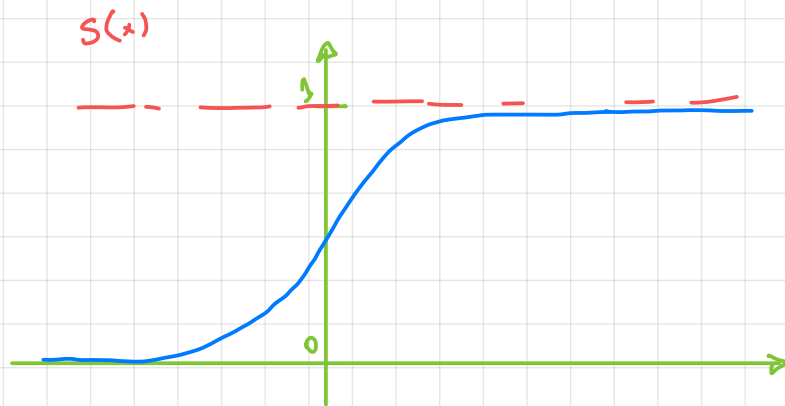
$$= \frac{1}{1 + e^{-\phi(\epsilon)}}$$

QED.

Ex: Logistic "activation" function

$$S(x) = \frac{1}{1 + e^{-cx}} \quad \text{for } c > 0.$$

$$\therefore S' = c \cdot S(1-S) > 0$$



$\therefore S$ is a "soft threshold"

Thrm: Probabilistic Modus Ponens

$$\text{Premise 1: } P(B|A) \geq c$$

$$\text{Premise 2: } P(A) \geq a$$

$$\text{Conclusion: } P(B) \geq a \cdot c$$

$$\text{Prf: } P(B) \geq P(A \cap B)$$

$$= P(A) \cdot \frac{P(A \cap B)}{P(A)}$$

$$= P(A) \cdot P(B|A)$$

$$\geq a \cdot c$$

monotonicity

$$P(A) > 0$$

defn cond. prob.

premise 1 and 2.

QED

$$\text{Check: } a = c = 1 \quad \therefore P(B) = 1$$

Thrm: Probabilistic Modus Tollens

$$\text{Premise 1: } P(B|A) \geq c > 0$$

$$\text{Premise 2: } P(B) \leq b.$$

$$\text{Conclusion: } P(A) \leq \min\left(1, \frac{b}{c}\right)$$

$$\text{Prf: } P(A) \leq P(A) \left(\frac{P(B)}{P(A \cap B)} \right) \quad \text{since } P(A \cap B) = P(B)$$

$$= \frac{P(B)}{P(B|A)}$$

$$\leq \frac{b}{P(B|A)}$$

premise 1.

$$\leq \frac{b}{c}$$

premise 2

$$\text{Always } P(A) \leq 1 \quad \therefore P(A) \leq \min\left(1, \frac{b}{c}\right)$$

QED.

Independent Sigma-Algebras (optional)

Defn: Sigma-algebras \mathcal{A} and \mathcal{B} on (Ω, \mathcal{A}, P) are independent iff $P(A \cap B) = P(A) \cdot P(B)$ for all $A \in \mathcal{A}$ and all $B \in \mathcal{B}$.

Thm: Events A and B are independent on probability space (X, \mathcal{A}, P) iff $\sigma(\{A\})$ and $\sigma(\{B\})$ are independent sigma-algebras

Prf: $\sigma(\{A\}) \stackrel{\text{CIA}}{=} \{\emptyset, A, A^c, X\}$ $\sigma(\{B\}) = \{\emptyset, B, B^c, X\}$.

"
←"

Say $\sigma(\{A\})$ and $\sigma(\{B\})$ are independent.

$\therefore A$ and B are independent because $A \in \sigma(\{A\})$ and $B \in \sigma(\{B\})$. QED.

"
→"

Say A and B are independent on (Ω, \mathcal{A}, P)

$$\therefore P(A \cap B) = P(A) \cdot P(B)$$

- $P(\Omega) = 1 = P(\Omega \cap \Omega) \quad \therefore P(\Omega \cap \Omega) = 1 = 1 \cdot 1 = P(\Omega) \cdot P(\Omega)$

$\therefore \Omega \in \sigma(\{A\})$ and $\Omega \in \sigma(\{B\})$ are independent.

- $0 = P(\emptyset) = P(\emptyset \cap \emptyset) \quad \therefore P(\emptyset \cap \emptyset) = 0 = 0 \cdot 0 = P(\emptyset) \cdot P(\emptyset)$.

- $P(A \cap \Omega) = P(A) = P(A) \cdot 1 = P(A) \cdot P(\Omega)$

$\therefore A \in \sigma(\{A\})$ and $\Omega \in \sigma(\{B\})$ are independent.

and same for $\Omega \in \sigma(\{A\})$ and $B \in \sigma(\{B\})$

$$\begin{aligned}
 - P(A^c \cap B^c) &\stackrel{\text{DeM}}{=} P((A \cup B)^c) = 1 - P(A \cup B) \\
 &\stackrel{\text{Add}}{=} 1 - P(A) - P(B) + P(A \cap B) \\
 &\stackrel{\text{Ind}}{=} 1 - P(A) - P(B) + P(A) \cdot P(B) \\
 &= (1 - P(A))(1 - P(B)) \\
 &= P(A^c) \cdot P(B^c)
 \end{aligned}$$

$$\therefore A^c \in \sigma(\{A\}) \text{ and } B^c \in \sigma(\{B\})$$

$$- A = A \cap \Omega = A \cap (B \cup B^c) \stackrel{\text{dist}}{=} (A \cap B) \cup (A \cap B^c)$$

$$\text{and } (A \cap B) \cap (A \cap B^c) = A \cap (B \cap B^c) = A \cap \emptyset = \emptyset.$$

$$\therefore P(A) \stackrel{\text{CAT}}{=} P(A \cap B) + P(A \cap B^c)$$

$$\therefore P(A \cap B^c) = P(A) - P(A \cap B)$$

$$\stackrel{\text{ind}}{=} P(A) - P(A) \cdot P(B)$$

$$= P(A)(1 - P(B))$$

$$= P(A) \cdot P(B^c).$$

$$\therefore A \in \sigma(\{A\}) \text{ and } B^c \in \sigma(\{B\}) \text{ are independent.}$$

Same argument, $A^c \in \sigma(\{A\})$ and $B \in \sigma(\{B\})$ are independent.

$$\therefore P(C \cap D) = P(C) \cdot P(D) \quad \forall C \in \sigma(\{A\}), \forall D \in \sigma(\{B\})$$

$$\therefore \sigma(\{A\}) \text{ and } \sigma(\{B\}) \text{ are independent.}$$

QED.

Defn: A π -system is a non-empty set collection closed under finite intersections

Thm: If probability measures P and Q agree on a π -system I (i.e. $P(A) = Q(A) \forall A \in I$) then P and Q agree on the generated sigma-algebra $\sigma(I)$.

Analogy: for functions continuous on $[0, T]$ you can uniquely represent f from samples $f(x)$ on rationals in the interval (or any set "dense" in $[0, T]$)

So the above theorem acts like a sampling theorem for probability measures on (X, \mathcal{Q})